

An Introduction to Functional Reactive Programming

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- Reactive Program: continually interacts with its environment in a timely manner.
- Examples: video games, mp3 players, robot controllers, aeroplane control systems . . .
- Contrast with
 - Interactive Programs, e.g. accessing a database
 - Transformational Programs, e.g. a compiler

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- FRP languages are domain-specific languages (the domain being reactive programming)
- Key characteristic: inherent notion of time
- Usually embedded in a host language (often Haskell)
- Also useful for modelling and simulation

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- Two choices:
 - continuous time: Time $\approx \{t \ge 0 \mid t \in \mathbb{R}\}$
 - discrete time: Time $\approx \mathbb{N}$

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- Two choices:
 - continuous time: Time $\approx \{t \ge 0 \mid t \in \mathbb{R}\}$
 - discrete time: Time $\approx \mathbb{N}$
- The original idea of FRP was to provide a continuous-time abstraction to the FRP programmer...

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- Two choices:
 - continuous time: Time $\approx \{t \ge 0 \mid t \in \mathbb{R}\}$
 - discrete time: Time $\approx \, \mathbb{N}$
- The original idea of FRP was to provide a continuous-time abstraction to the FRP programmer...
- ... while automating the discretisation necessary for implementation.

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Signal a \approx Time \rightarrow a

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Signal a \approx Time \rightarrow a

• There are also instantaneous occurrences called events.

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- There are also instantaneous occurrences called events.
- In a discrete-time setting, events can be embedded within signals:

Event a = Signal (Maybe a)

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Signal a \approx Time \rightarrow a

- There are also instantaneous occurrences called events.
- In a discrete-time setting, events can be embedded within signals:

Event a = Signal (Maybe a)

• In a continuous-time setting, they require a separate abstraction:

Event $a \approx [(Time, a)]$

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- FRP languages keep signals abstract, providing several signals, and functions on signals, as primitives.
- Some go further and only provide functions on signals as a first-class abstraction.
- These are called signal functions:

SF a b pprox Signal a ightarrow Signal b

- The abstraction prevents many "bad" signal functions from being defined.
- E.g. causality can be enforced (the present cannot depend on the future).

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- A DSL embedded in Haskell
- No signals, only signal functions
- Pretends to have continuous time
- Has been used for a variety of applications: video games, sound synthesis, robot simulators, GUIs, virtual reality, visual tracking, animal monitoring...

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Arrows

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Implementation

Summary

Some Yampa Primitives

Example Primitives

 $\begin{array}{l} \mbox{constant :: } b \ \rightarrow \ SF \ a \ b \\ \mbox{integral :: } Num \ a \ \Rightarrow \ SF \ a \ a \\ \mbox{delay :: } Time \ \rightarrow \ a \ \rightarrow \ SF \ a \ a \\ \mbox{edge :: } SF \ Bool \ (Event \ ()) \\ \mbox{tag :: } Event \ a \ \rightarrow \ b \ \rightarrow \ Event \ b \\ \mbox{switch :: } SF \ a \ (b, \ Event \ e) \ \rightarrow \ (e \ \rightarrow \ SF \ a \ b) \ \rightarrow \ SF \ a \ b \end{array}$

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The Arrow Framework

Arrow Classes

```
class Category (a :: * \to * \to *) where

id :: a b b

(\circ) :: a c d \to a b c \to a c d

class Category a \Rightarrow Arrow (a :: * \to * \to *) where

arr :: (b \to c) \to a b c

first :: a b c \to a (b, d) (c, d)

class Arrow a \Rightarrow ArrowLoop (a :: * \to * \to *) where

loop :: a (b, d) (c, d) \to a b c
```

Yampa Signal Functions are Arrows

```
instance Category SF where ...
```

```
instance Arrow SF where ...
```

```
instance ArrowLoop SF where ...
```

Arrows

Examples

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Signal Functions Combinators







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$$\begin{array}{ll} \operatorname{arr} & :: (a \to b) \to SF \ a \ b \\ (\ggg) & :: SF \ a \ b \to SF \ b \ c \to SF \ a \ c \\ \operatorname{return} A :: SF \ a \ a \\ \operatorname{first} & :: SF \ a \ c \to SF \ (a, b) \ (c, b) \\ \operatorname{second} & :: SF \ b \ c \to SF \ (a, b) \ (a, c) \\ (\&\&) & :: SF \ a \ b \to SF \ a \ c \to SF \ a \ (b, c) \\ \operatorname{loop} & :: SF \ (a, c) \ (b, c) \to SF \ a \ b \end{array}$$

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Example Signal Functions

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Example Signal Functions

• See accompanying code for bouncing-ball example...

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The SF data type (simplified)

data SF a $b \approx$ SF (DTime \rightarrow a \rightarrow (SF a b, b))

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The SF data type (simplified)

data SF a $b \approx$ SF (DTime \rightarrow a \rightarrow (SF a b, b))

An alternative implementation

data SF a $b \approx$ SF (DTime \rightarrow s \rightarrow a \rightarrow (s, b))

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- FRP languages usually take one of two implementation strategies:
 - Pull-driven implementations update at every time step (good for systems with continuously changing signals).
 - Push-driven implementations only update when an event occurs, and only the parts of the program that depend on that event (good for systems with signals that change at discrete points in time).
- In practice, FRP implementations contain a lot of optimisations to avoid unnecessary computation.
- But efficient implementation of FRP remains an open problem.

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- FRP languages are domain-specific languages for reactive programming.
- Their key characteristic is an implicit notion of time.
- If you want to experiment with Yampa, I'd recommend Henrik Nilsson's recent mini-course:

http://www.cs.nott.ac.uk/~nhn/ITU-FRP2010/

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